

LISP: A Level of Indirection for Routing

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RIPE 55

Amsterdam



- What is LISP?
- Why LISP?
- Protocol Details
 - Data-plane part
 - Control-plane part
- Prototype & Deployment Status
- Summary

What is LISP?

- Locator/ID Separation Protocol
 - Details to follow...
- Design Space:
 - Scalability of the routing/addressing system
- Design Assumptions/Goals:
 - Network-based solution
 - No changes to hosts whatsoever
 - No new addressing changes to site devices
 - Very little configuration file changes
 - Must be incrementally deployable
 - Address family agnostic

Locator/ID Separation?

- The idea here is that the IP address is overloaded
 - It encodes both location in the topology (locator) and the identity of the user of the address
- The locator role is used by the routing system
- The idenity role is used by upper layer protocols
 - e.g., TCP psuedo-header
- Since we want locators to aggregate topologically, and since identities are usually allocated on organizational boundaries, it is difficult to get one number space to efficiently serve both purposes
- One solution is to split the functions -- This is at the heart of the Locator/ID split idea
 - So how might we achieve this?

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Implementing a Locator/ID Split

- There are two main ways to implement the Loc/ID split
- Address Rewriting (aka "Network NAT")
 - If you have enough address space (e.g., IPv6), you could use the lower 64 bits as an identifier, and the upper 64 bits as a locator, and rewrite the locator at the border
 - This is the basis of O'Dell's 8+8/GSE scheme

Map-n-Encap

- You could also put another header on the packet, and make the inner header carry the IDs and the outer header carry the locators
- LISP is an instance of this approach

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So What is LISP?

 LISP separates out location and identification from an existing IP address semantic



Why the Separation?

- The level of indirection allows us to:
 - Keep either ID or Location fixed while changing the other
 - Creates separate namespaces which can have different allocation properties
- In particular, the EID allocation hierarchy can follow a different topology than the RLOC allocation hierarchy
 - Which must at least be somewhat congruent to network topology if we want to be able to aggregate effectively (Rekhter's Law)

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Why the Separation?

- By keeping IDs fixed
 - Assign fixed addresses that never change to hosts and routers at a site
- You can change Locators
 - Now the sites can change providers
 - Now the hosts can move

Why LISP?

- Operationally
 - Improve site multihoming
 - Improve ISP Traffic Engineering
 - Reduce site renumbering costs
 - Reduce size of core routing tables
 - Conserve IPv4 (and IPv6) address space
 - PI for all?
 - Some form of mobility?
- Architecturally
 - Create two namespaces: EIDs and Locators

What Provoked This?

- Stimulated from problem statement effort at the Amsterdam IAB Routing Workshop on October 18/19 2006
 RFC 4984
- More info on problem statement:
 - http://www.vaf.net/~vaf/apricot plenary.pdf

- Data-plane
 - Design for encapsulation and tunnel router placement
 - Design for locator reachability
 - Data triggered mapping service
- Control-plane
 - Design for a scalable mapping service
 - Examples are: CONS, NERD, and RPMD

Jack-Up Model



LISP is a Jack-Up Model

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LISP Packet Format

٥ 1 З 67890123456789012345678901 0 1 2 3 4 5 -+-+-+-+-+-+-+-+-+ Version IHL Type of Service Total Length Identification Flags Fragment Offset 1 OH Time to Live | Protocol = 17 | Header Checksum \ Source Routing Locator Destination Routing Locator Source Port Dest Port \ UDP length UDP Checksum Type Locator Reach Bits Nonce Nonce / Version IHL Type of Service Total Length Flags Identification Fragment Offset 1 тн Time to Live Protocol Header Checksum \ Source EID Destination EID

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- When there is no mapping in the ITR:
 - Use data-triggered UDP Map-Reply
 - Invoking by sending outer DA to inner DA
 - Send on alternative topology
 - BGP-over-GRE using EIDs as NLRI
 - No mapping data in BGP
 - No changes to BGP

- When there is no mapping in the ITR:
 - ITR sends CONS Map-Request
 - ITR gets Map-Reply
 - Packets get dropped in the meantime
 - Only happens first time when source site talks to destination site
 - Scalable because EID-prefix allocation not tied to underlying topology
 - Pull model

- When there is no mapping in the ITR:
 - Lets have mappings always in ITRs
 - NERD pushes a signed file
 - RPMD pushes signed records
 - ITRs never table-miss at expense of compressed data-set sent to every ITR
 - Push model

Prototype

- cisco has a LISP prototype implementation
 - Started the week of IETF Prague (March 2007)
- OS platform is DC-OS
 - Linux underlying OS
- Hardware platforrm is Titanium
 - 1 RU dual-core off-the-shelf PC with 7 GEs
- Based on draft-farinacci-lisp-04.txt
- Software switching only
- Supports both IPv4 and IPv6

Prototype

- Supports ITR encap and ETR decap
 - Load-balancing among locators
 - Respects priority & weight per mapping
- Multiple EID-prefixes per site
- Support for locator reachability
- Multi-VRF support for BGP-over-GRE
- Supports both IPv4 and IPv6

Prototype - What's Next?

- Implement "crossover" support
 - IPv6-EIDs over IPv4-Locators
 - IPv4-EIDs over IPv6-Locators
- Implement shortest-path Mobility
 - Use route-returnability check to protect ITR spoofing
- Start CONS implementation

Prototype Testing

- Dino's Unit Testing
- Meyer, Fuller, Lewis, Shepherd testing since July 2007
- External Pilot
 - Shooting for post Vancouver IETF
 December 2007

Internal Pilot Testing



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Summary

- LISP: draft-farinacci-lisp-04.txt
- CONS: draft-meyer-cons-02.txt
- NERD: draft-lear-nerd-02.txt
- Please send us your comments!
- Please let us know if you're interested in pilot deployment

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Questions/Comments?

Thanks!